



SCION: A Secure Multipath Interdomain Routing Architecture

Adrian Perrig

Network Security Group, ETH Zürich

ETH zürich

SCION

Internet Weakness: DoS and DDoS Attacks

- Expensive and difficult to protect against DoS und DDoS attacks
- Despite large investments, attacks continue to be successful
 - **November 2015: Protonmail attacked during 1 week**
 - **March 2016: CH e-commerce under attack: Digitec, Galaxus, SBB, Migros, etc.** (Hackers demanded 25 Bitcoins to stop attacks)
 - **Fall 2016: Global Mirai botnet attacks, e.g., OVH, Dyn, russian banks**
 - June 2017: Northkorea “Hidden Cobra” botnet uncovered
 - September 2017: Global airport chaos, DDoS paralyzes checkin systems
- **Can we reliably defend against DDoS attacks?**

Internet Weakness: Communication Path Hijacking

- Sender und receiver have limited control over routing paths
- Attacks can hijack and relay paths
- **How can we guarantee communication paths?**



Internet Weakness: Kill Switch ruptures Sovereignty

- Current Internet suffers from several “Kill Switches”, which can halt communication within a geographical area
- Several attack avenues exist: DDoS, BGP hijacking, DNS redirection, BGPSEC / DNSSEC / TLS certificate revocation
- Example August 2017: An erroneous route injected by Google prevents communication for 50% of Internet in Japan during 40 minutes
- **Can we construct an Internet without Kill Switches?**

SCION Architecture Design Goals

- **High availability**, even for networks with malicious parties
 - Adversary: access to management plane of router
 - Communication should be available if adversary-free path exists
- **Secure entity authentication**
that scales to global heterogeneous (dis)trusted environment
- **Flexible trust**: enable selection of trust roots
- **Transparent operation**: clear what is happening to packets and whom needs to be relied upon for operation
- **Balanced control** among ISPs, senders, and receivers
- **Scalability, efficiency, flexibility**

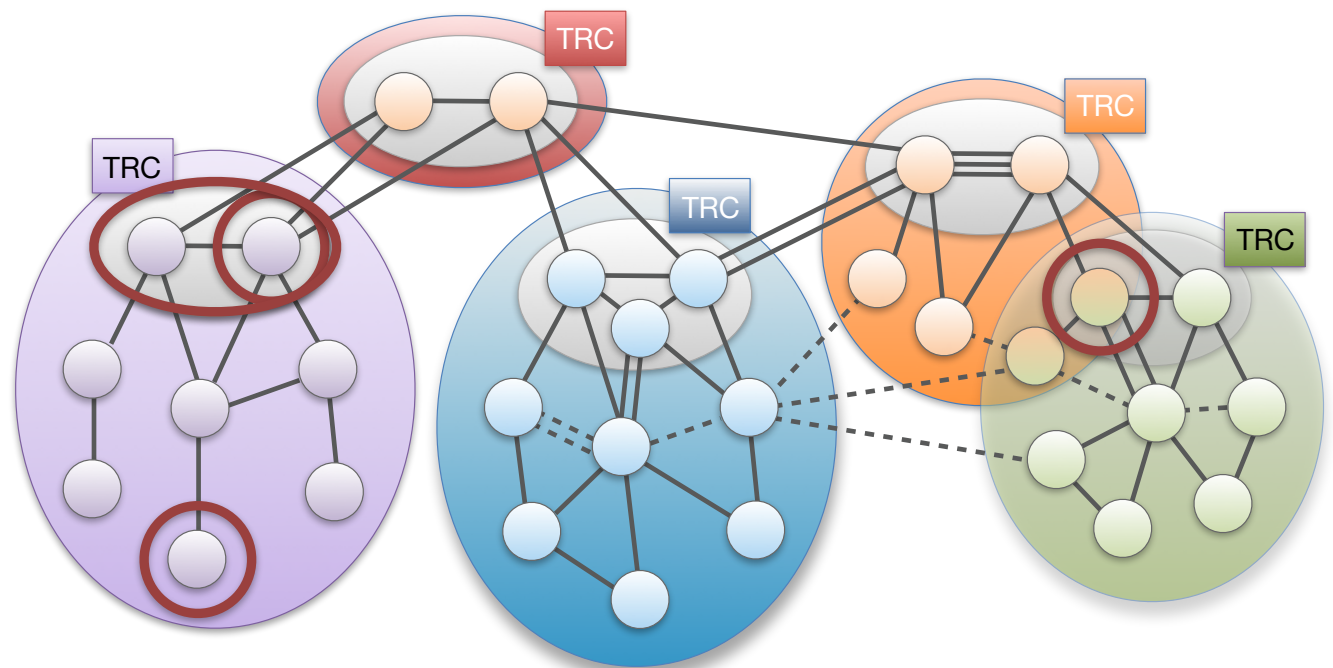


SCION Overview

- Control plane: How to find end-to-end paths?
 - Path exploration
 - Path registration
- Data plane: How to send packets
 - Path lookup
 - Path combination
- Deployment
- Demos

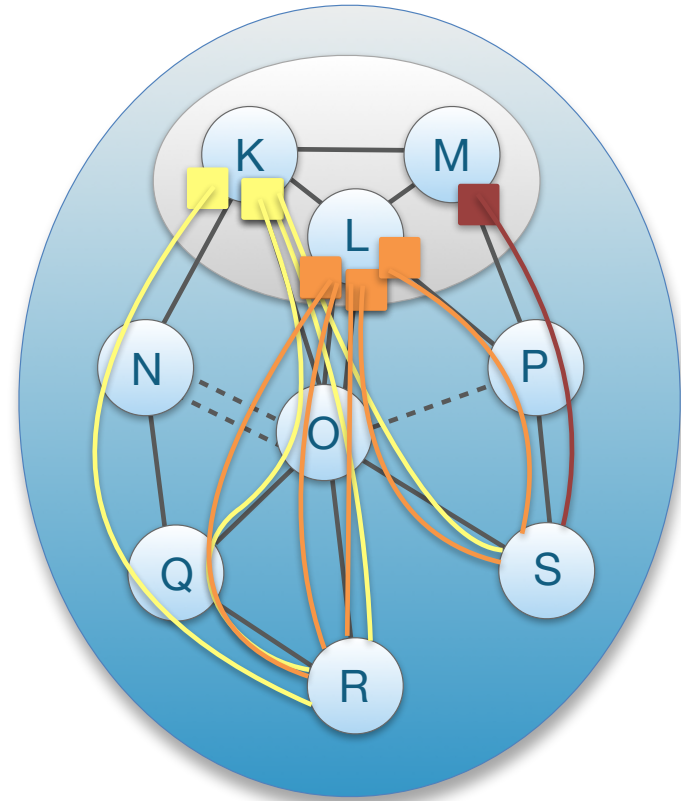
Approach for Scalability: Isolation Domain (ISD)

- Isolation Domain (ISD): grouping of ASes
- ISD core: ASes that manage the ISD
- Core AS: AS that is part of ISD core
- Control plane is organized hierarchically
 - Inter-ISD control plane
 - Intra-ISD control plane



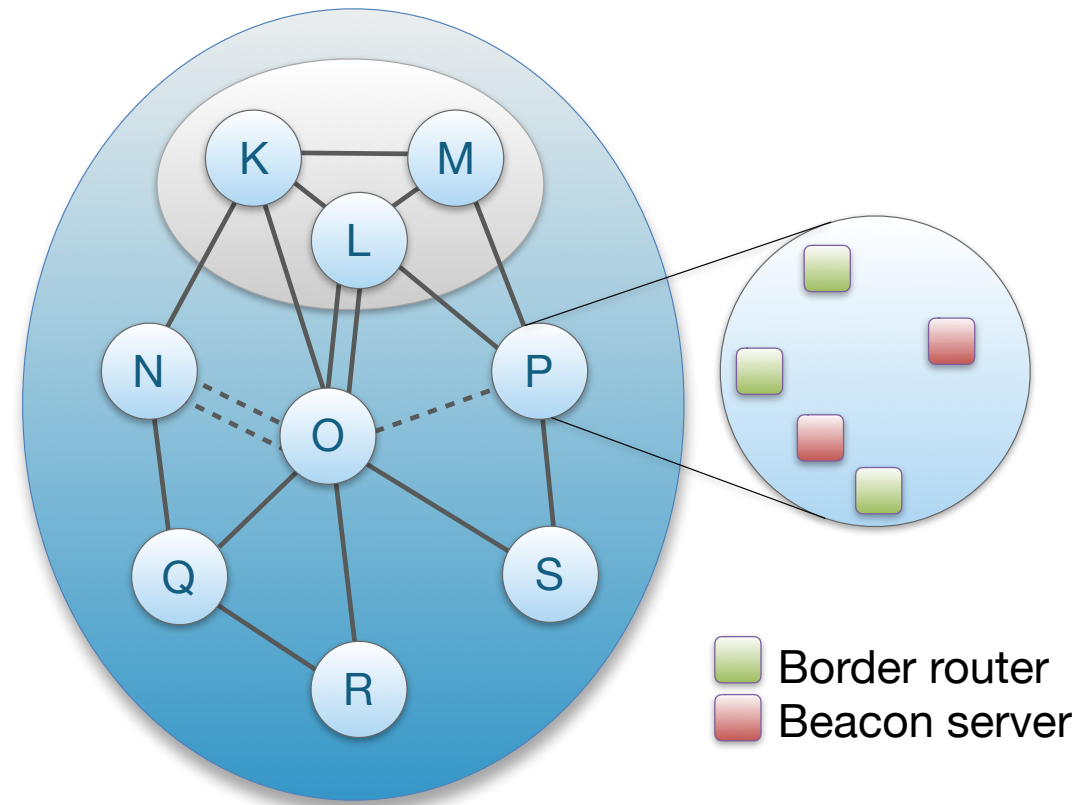
Intra-ISD Path Exploration: Beaconing

- Core ASes K, L, M initiate Path-segment Construction Beacons (PCBs), or “beacons”
- PCBs traverse ISD as a flood to reach downstream ASes
- Each AS receives multiple PCBs representing path segments to a core AS



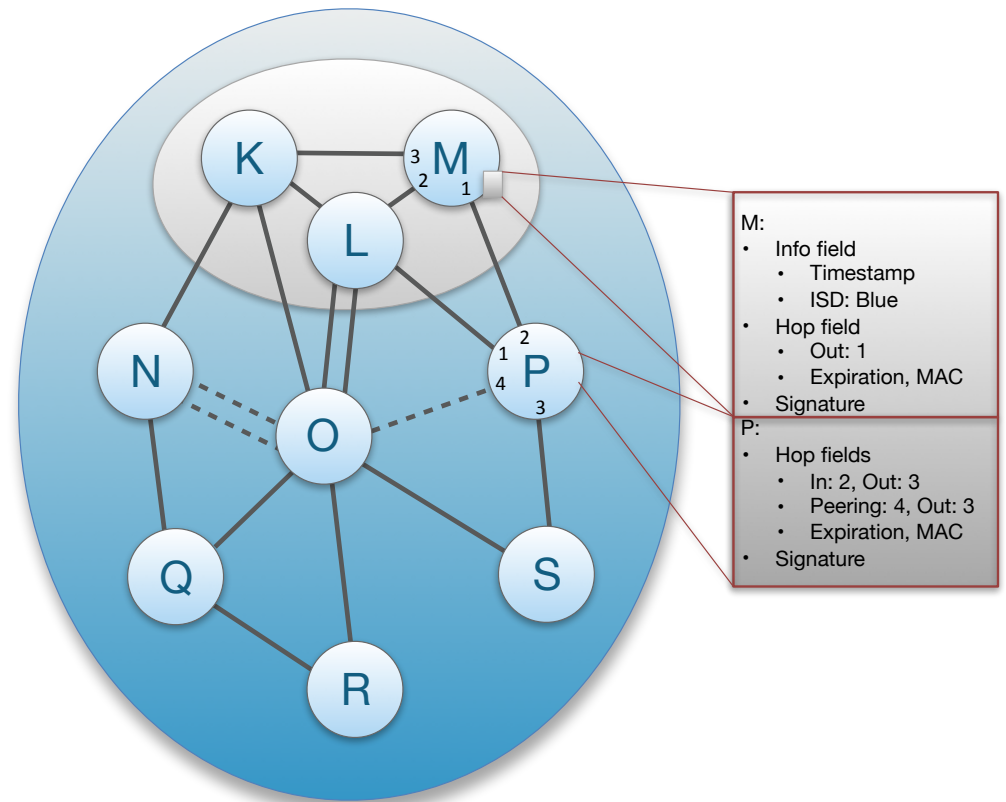
Beaconing in More Detail

- Each AS deploys one or multiple beacon servers
- PCBs are sent via a SCION service anycast packet
- SCION border routers receive PCB and select one beacon server to forward it to
- Beacon servers coordinate to re-send PCBs periodically to downstream ASes
 - Currently every 5 seconds, PCBs are selected and forwarded



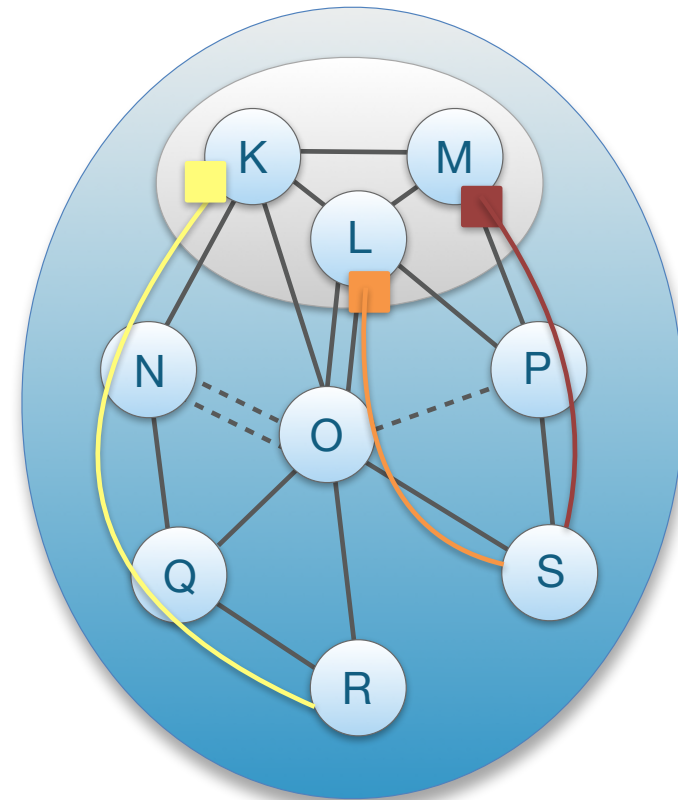
PCB Contents

- A PCB contains an info field with:
 - PCB creation time
- Each AS on path adds:
 - AS name
 - Hop field for data-plane forwarding
 - Link identifiers
 - Expiration time
 - Message Authentication Code (MAC)
 - AS signature

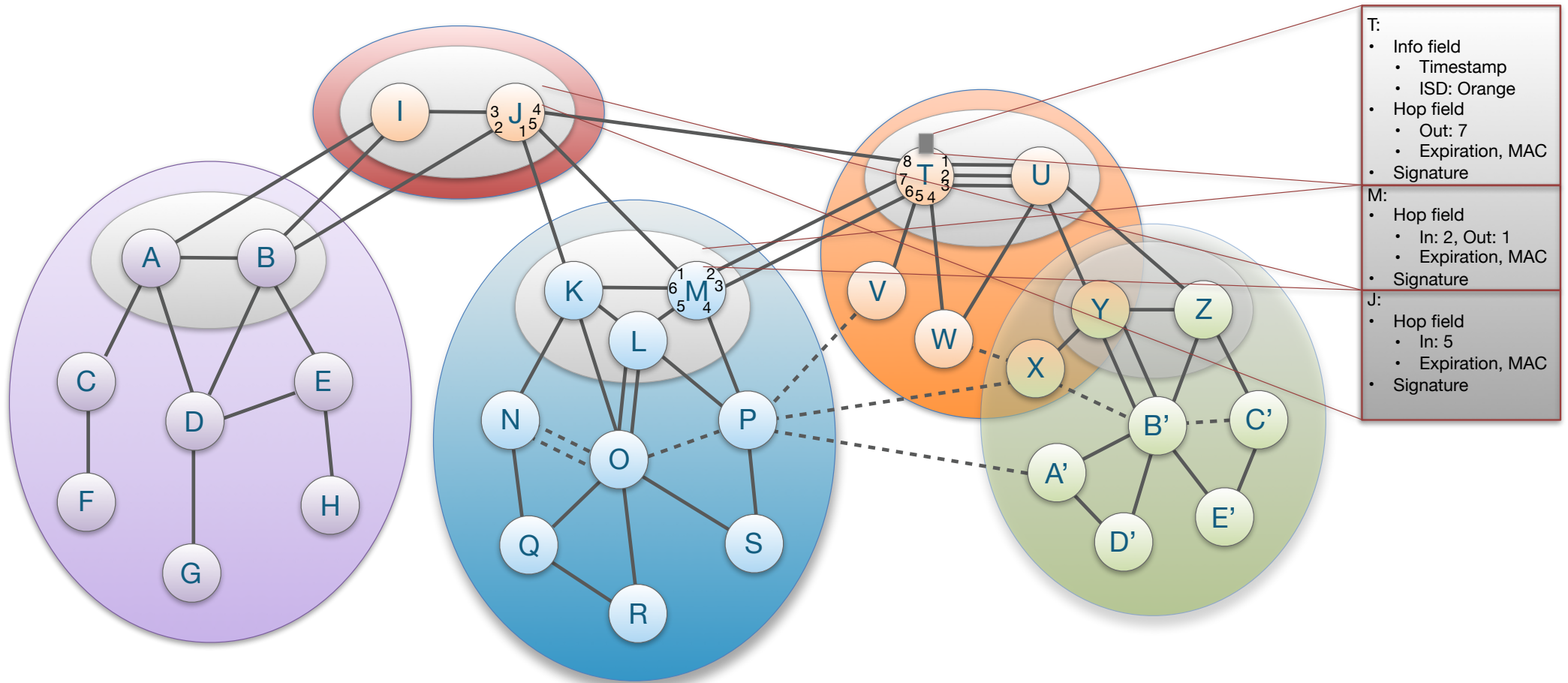


Up-Path and Down-Path Segments

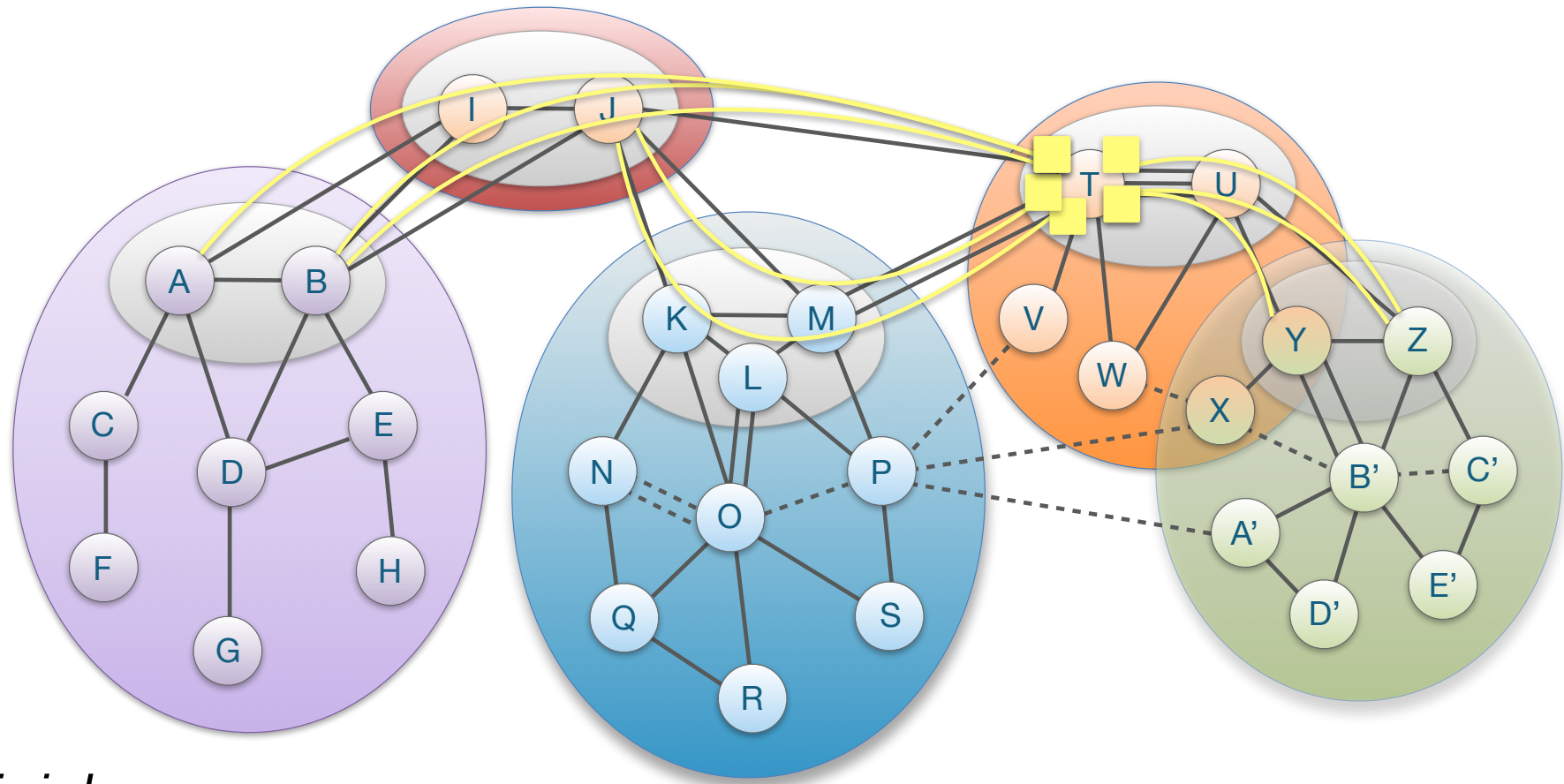
- Intra-ISD beaconing process sends PCBs to ASes
- PCBs contain **path segments** that can be used as communication paths to communicate with the core AS that initiated it
- **Up-path segment**: PCB is used from AS to core AS
 - Example: $R \rightarrow K$
- **Down-path segment**: PCB is used from core AS to AS
 - Example: $M \rightarrow S$



Core Beaconsing for Inter-ISD Path Exploration

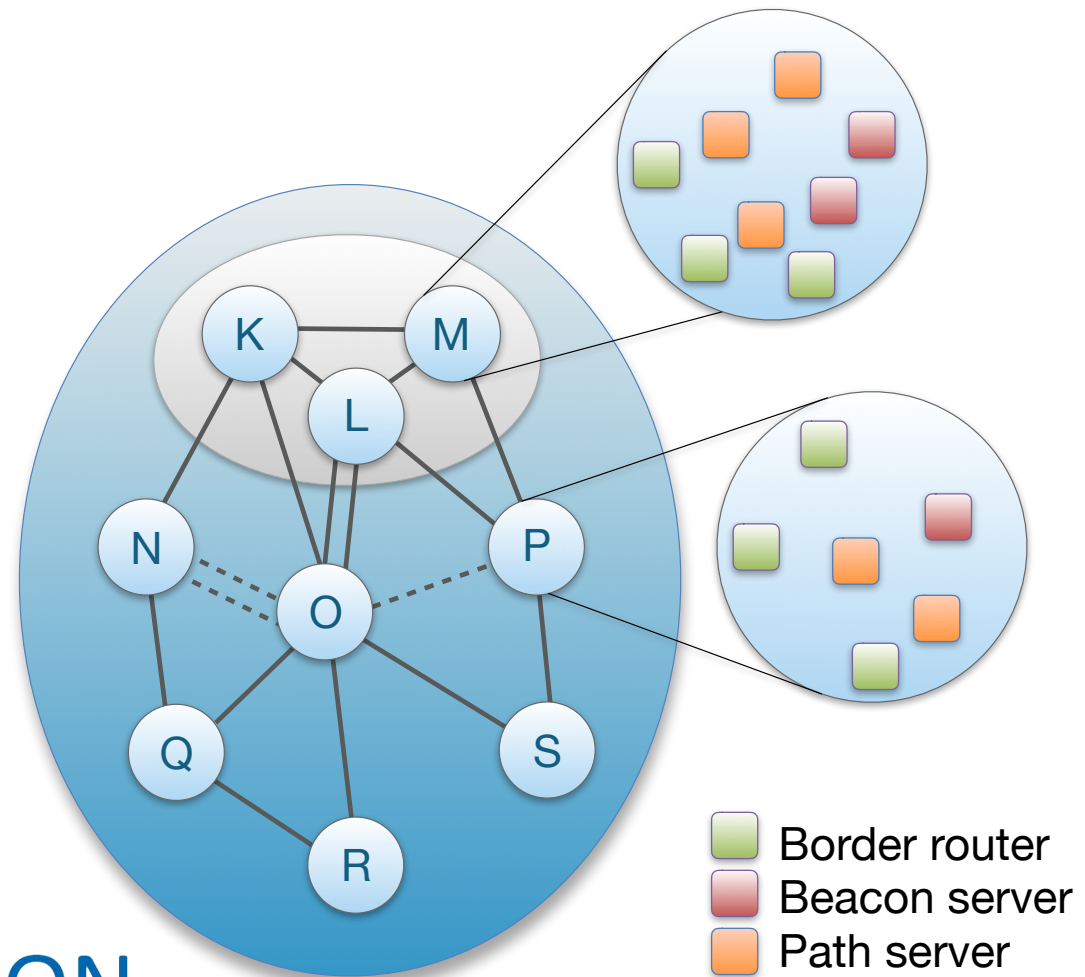


Inter-ISD Path Exploration: Sample Core-Path Segments from AS T



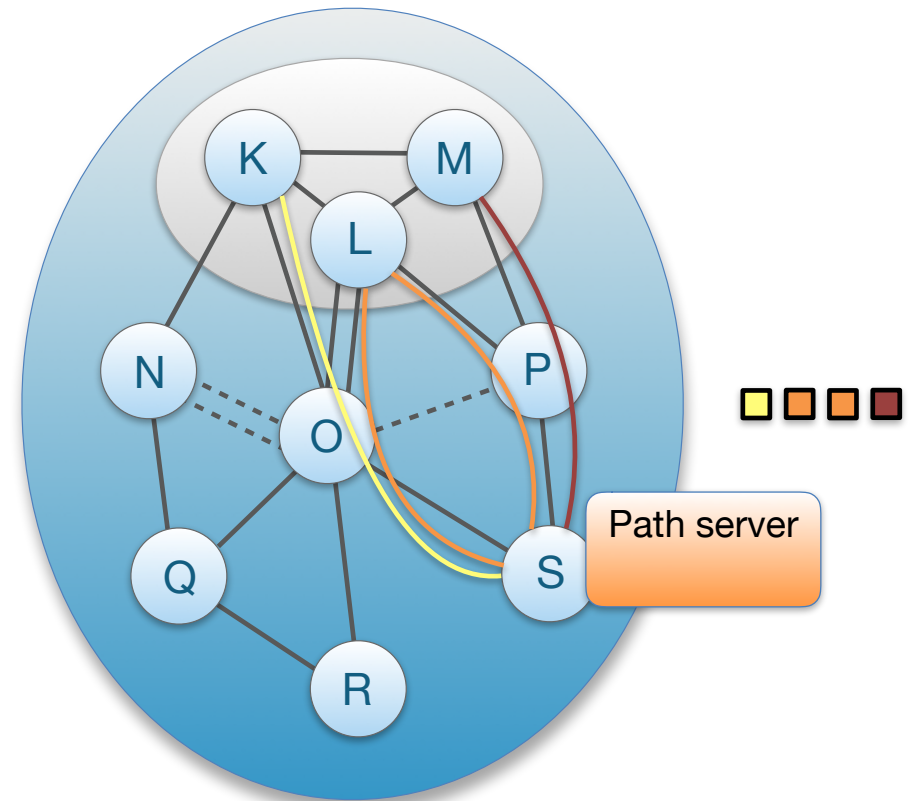
Path Server Infrastructure

- Path servers offer lookup service:
 - ISD, AS → down-path segments, core-path segments
 - Local up-path segment request → up-path segments to core ASes
- Core ASes operate core path server infrastructure
 - Consistent, replicated store of down-path segments and core-path segments
- Each non-core AS runs local path servers
 - Serves up-path segments to local clients
 - Resolves and caches response of remote AS lookups



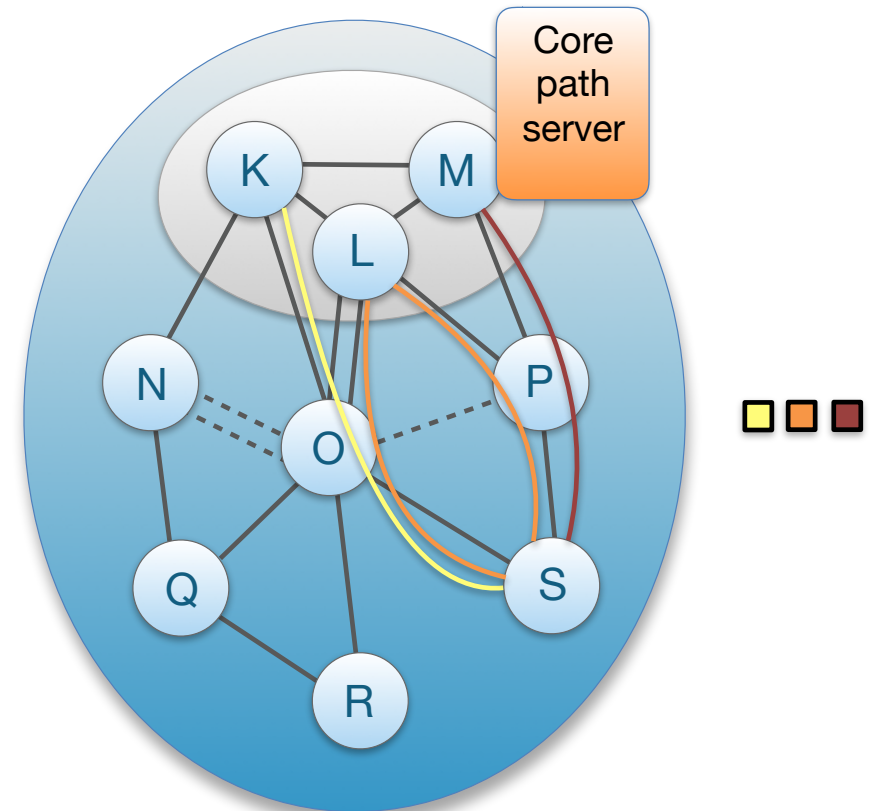
Up-Path Segment Registration

- AS selects path segments to announce as **up-path segments** for local hosts
- Up-path segments are registered at local path servers



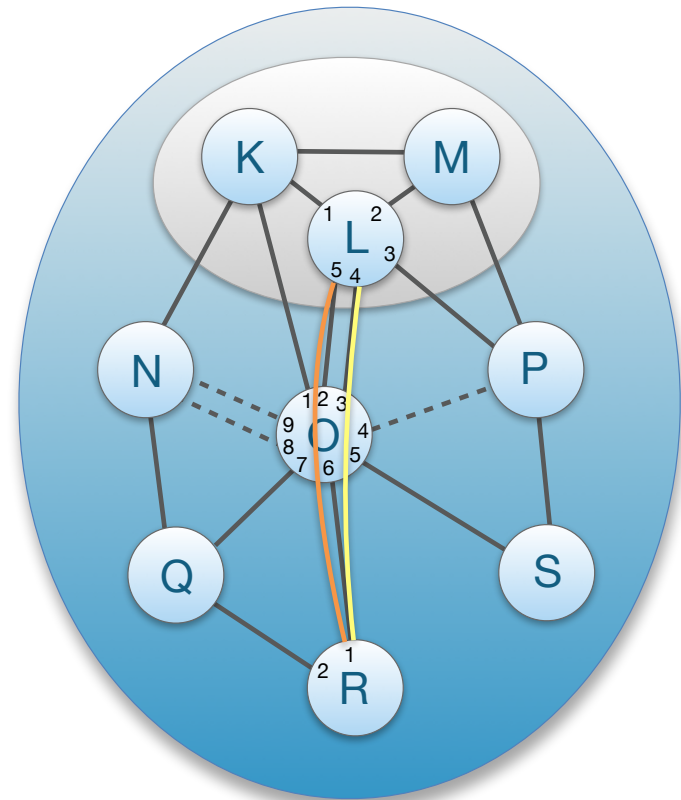
Down-Path Segment Registration

- AS selects path segments to announce as **down-path segments** for others to use to communicate with AS
- Down-path segments are uploaded to core path server in core AS



Ingress and Egress Interface Identifiers

- Each AS assigns a unique integer identifier to each interface that connects to a neighboring AS
- The interface identifiers identify ingress/egress links for traversing AS
- ASes use internal routing protocol to find route from ingress SCION border router to egress SCION border router
- Examples
 - Yellow path: L:4, O:3,6, R:1
 - Orange path: L:5, O:2,6, R:1



SCION Overview

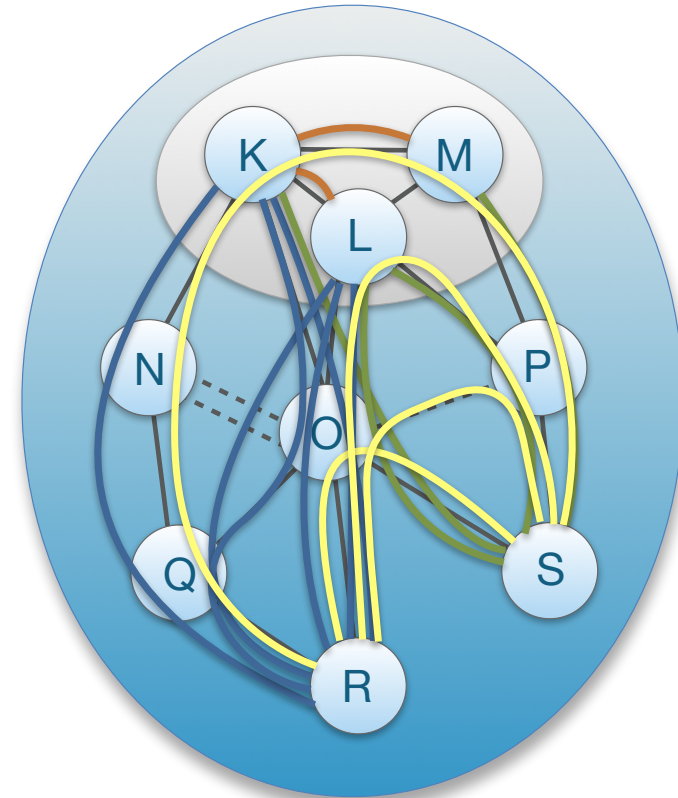
- Control plane: How to find end-to-end paths?
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Path Lookup

- Steps of a host to obtain path segments
 - Host contacts RAINS server with a name
H → RAINS: www.scion-architecture.net
RAINS → H: ISD X, AS Y, local address Z
 - Host contacts local path server to query path segments
H → PS: ISD X, AS Y
PS → H: up-path, core-path, down-path segments
 - Host combines path segments to obtain end-to-end paths, which are added to packets

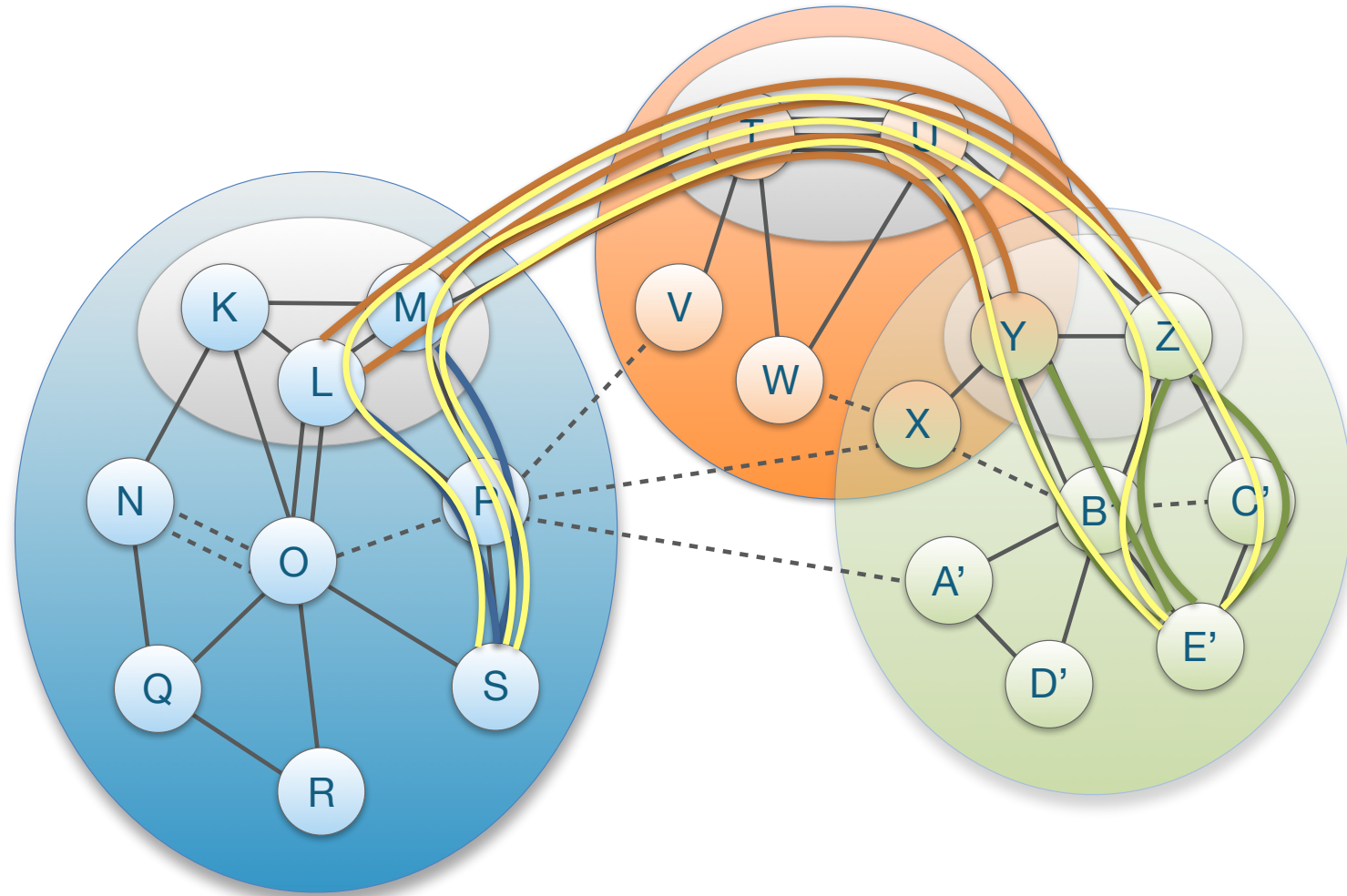
Path Lookup: Local ISD

- Client requests path segments to $\langle \text{ISD}, \text{AS} \rangle$ from local path server
- If down-path segments are not locally cached, local path server send request to core path server
- Local path server replies
 - Up-path segments to local ISD core ASes
 - Down-path segments to $\langle \text{ISD}, \text{AS} \rangle$
 - Core-path segments as needed to connect up-path and down-path segments

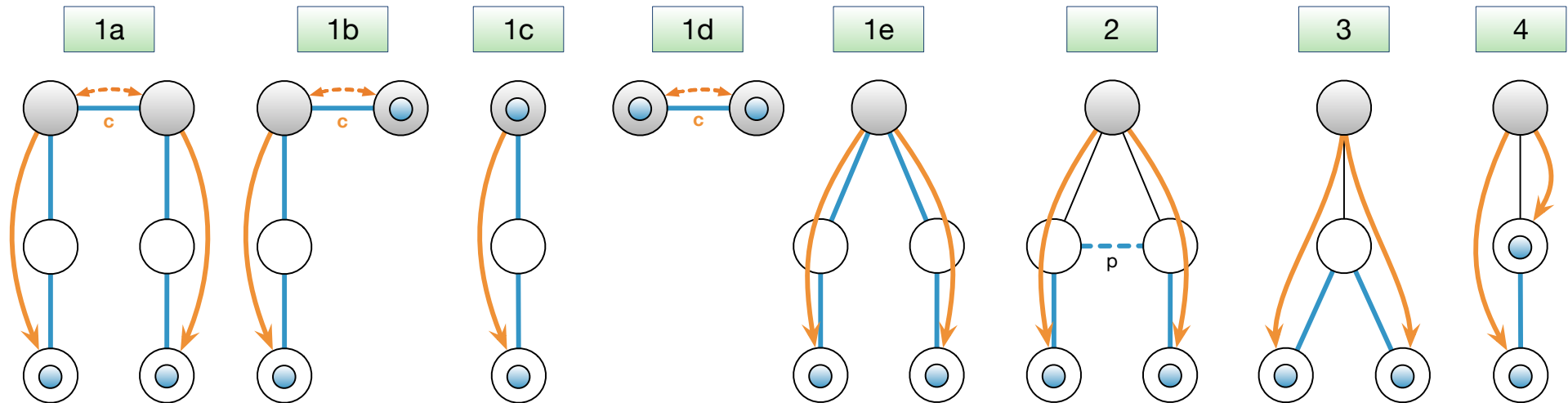


Path Lookup: Remote ISD

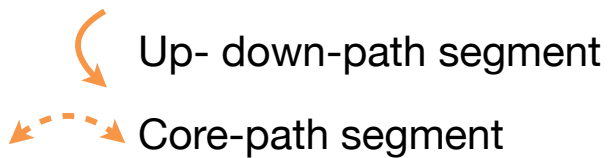
- Host contacts local path server requesting <ISD, AS>
- If path segments are not cached, local path server will contact core path server
- If core path server does not have path segments cached, it will contact remote core path server
- Finally, host receives up-, core-, and down-segments



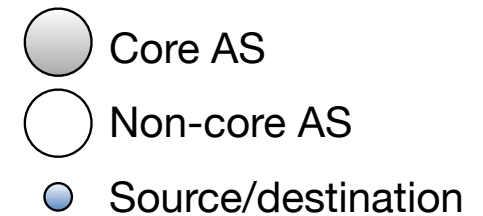
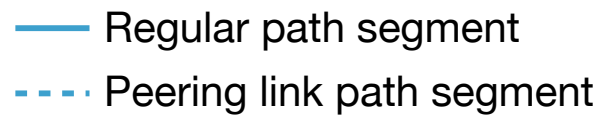
Path Combination



Control-plane path segments:

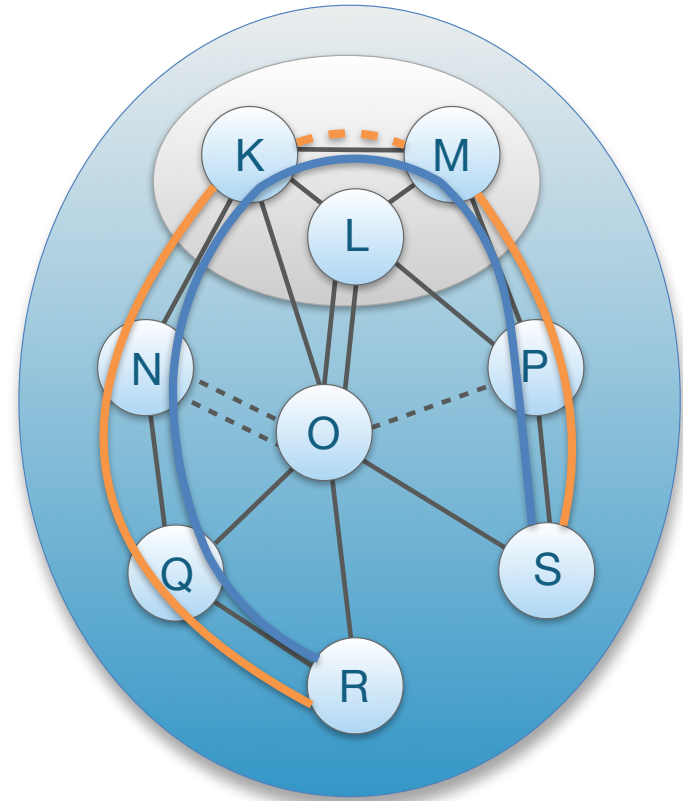
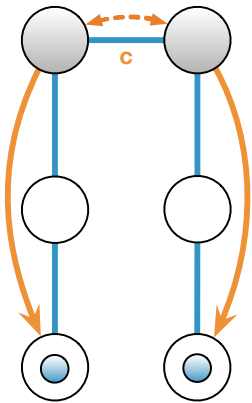


Data-plane paths:



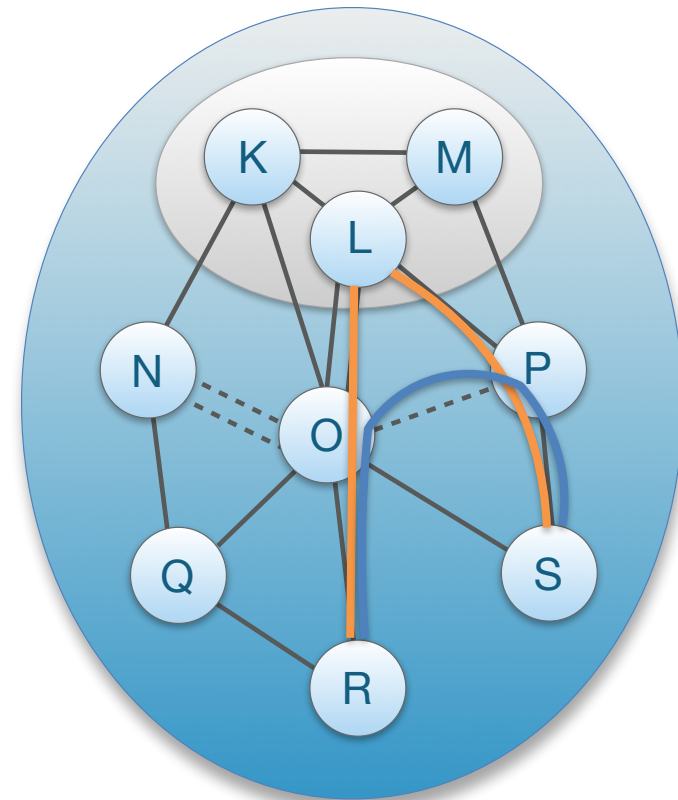
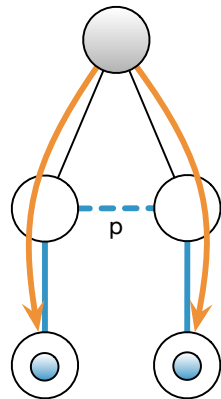
Path Combination Example (1)

- Core-segment combination:
Up-path segment +
core-path segment +
down-path segment



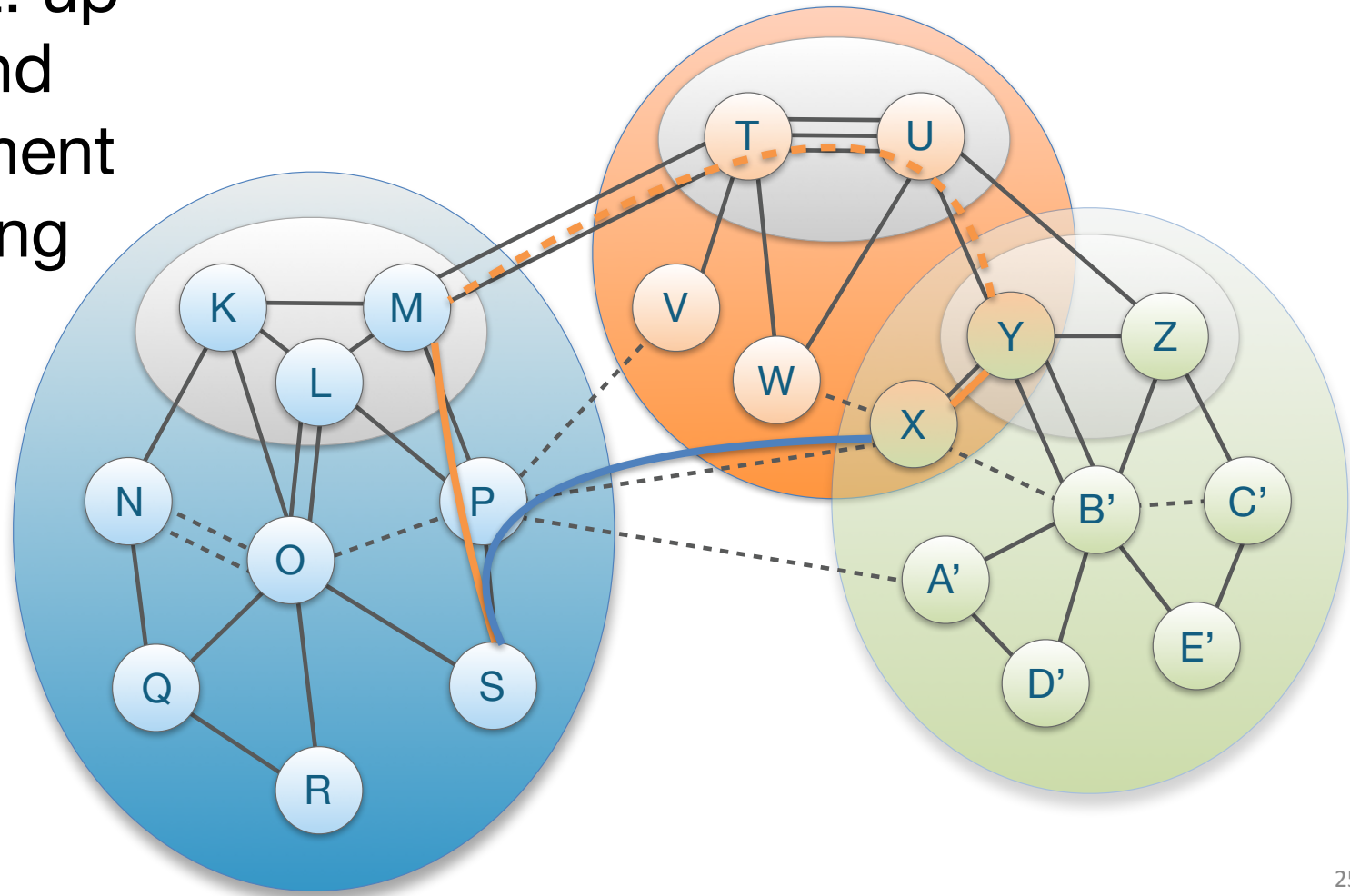
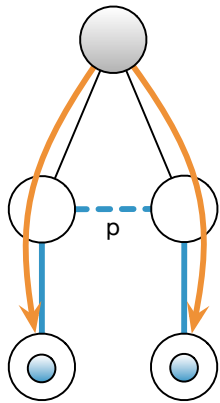
Path Combination Example (2)

- Peering shortcut: up-path segment and down-path segment offer same peering link



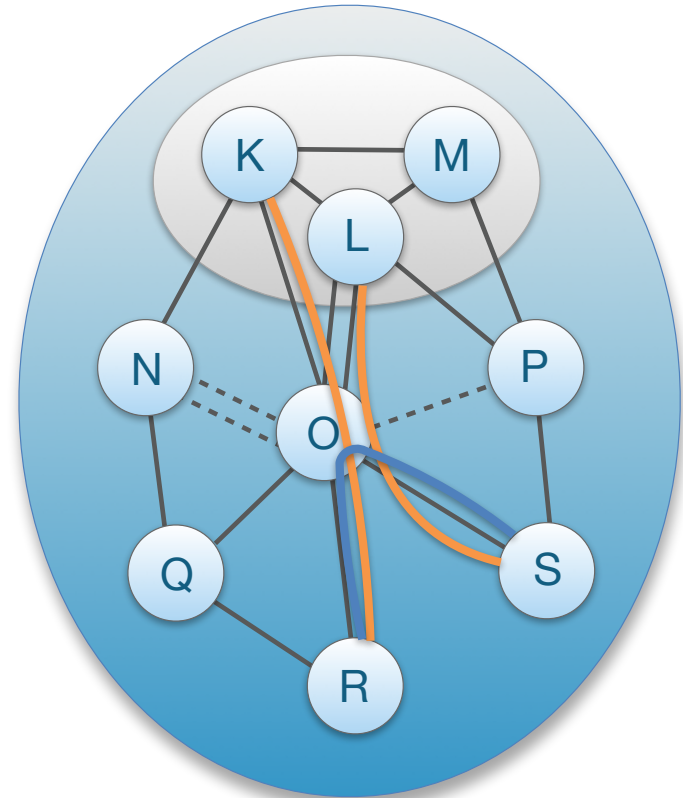
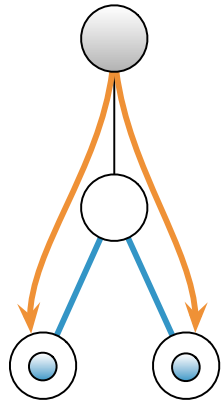
Path Combination Example (3)

- Peering shortcut: up-path segment and down-path segment offer same peering link

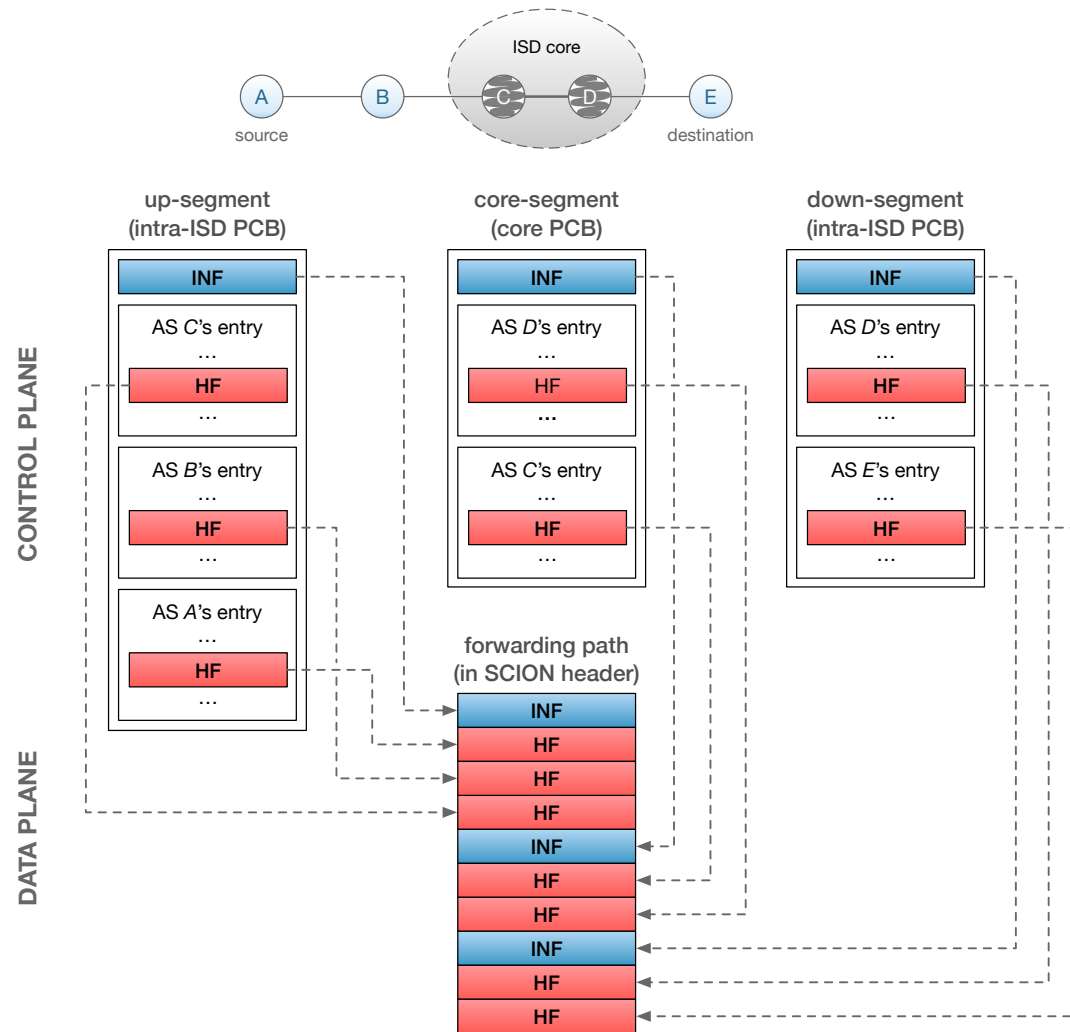


Path Combination Example (4)

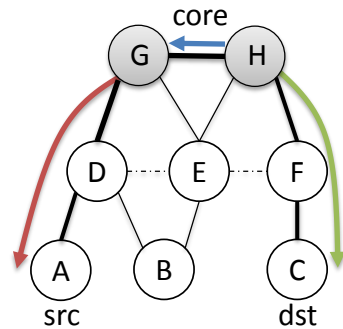
- AS shortcut path through common AS on up-path and down-path segment



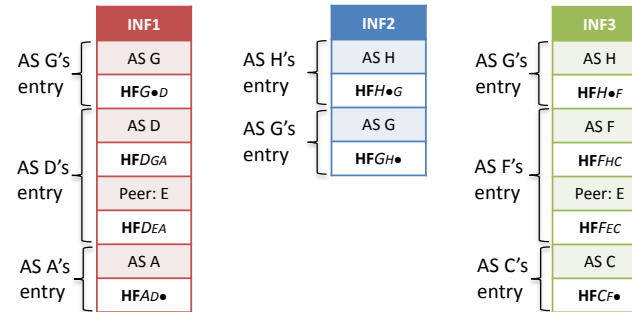
Path Construction



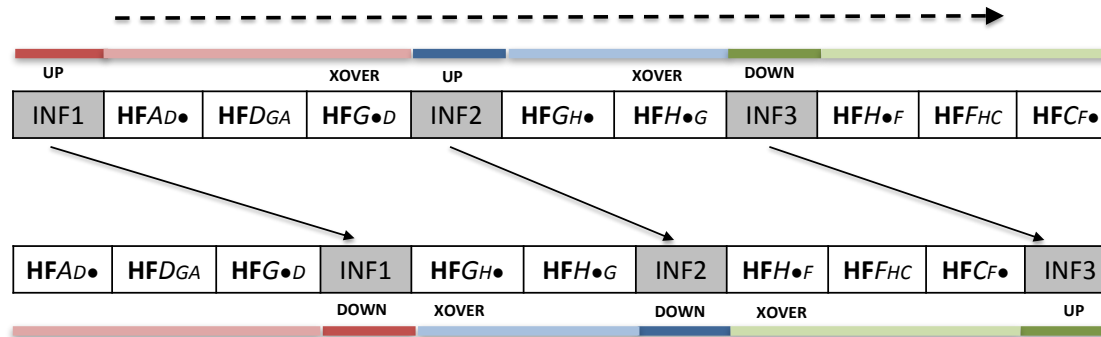
Path Encoding in Packet



Path segments:

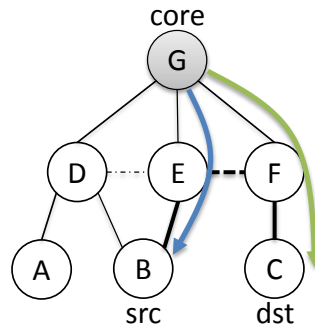


source to destination path

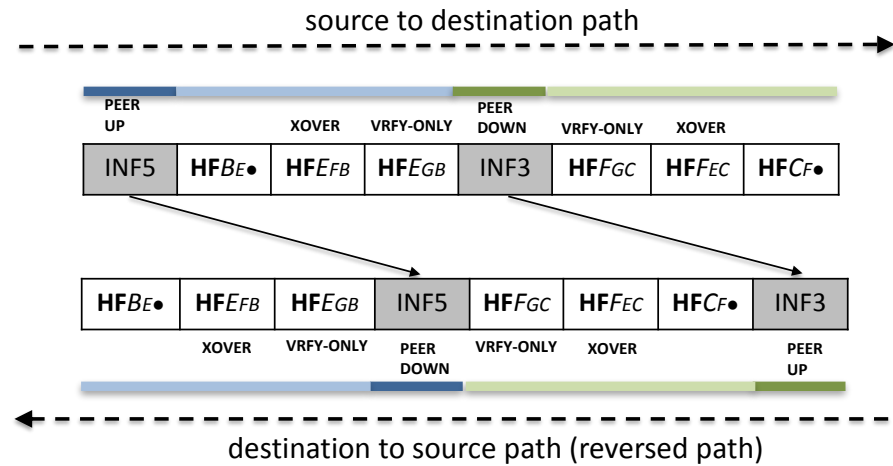
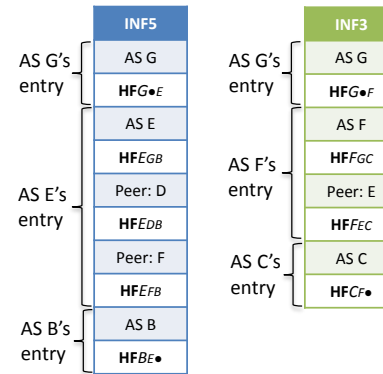


destination to source path (reversed path)

Path Encoding in Packet



Path segments:

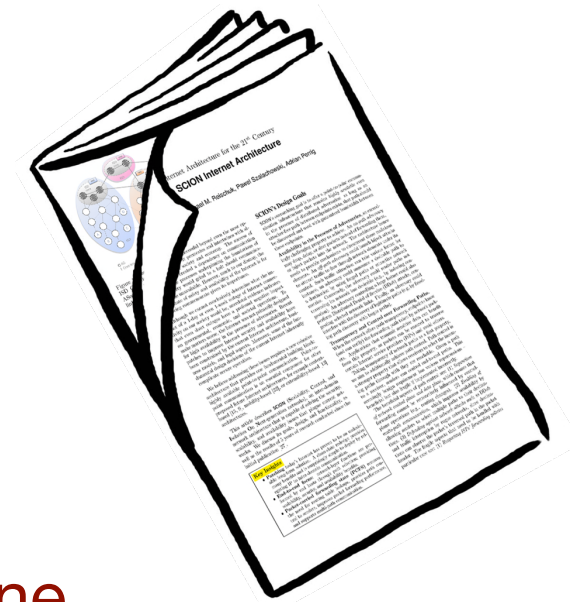


Hop Field MAC Verification

- Message Authentication Code (MAC) computation and verification of Hop Field MAC value based on local AS secret key
 - Key is not shared with any external entity
- Computation: $MAC_K(\text{Timestamp}, \text{Flags}'_{HF}, \text{ExpTime}, \text{Ingress}, \text{Egress}, \text{HF}')$
 - HF' is hop field of previous AS
- In most cases, HF' size is 8 bytes, so MAC computation can be done over 128 bits: with CMAC and AES, only a single encryption operation is needed
- With AESni HW crypto, only ~50 cycles are needed to compute MAC!
 - Note that a DRAM memory lookup takes ~200 cycles
 - AES operation requires less energy than TCAM lookup
 - Thus, **SCION forwarding can be faster and require less energy than IP forwarding**

SCION Summary

- Complete re-design of network architecture resolves numerous fundamental problems
 - BGP protocol convergence issues
 - Separation of control and data planes
 - Isolation of mutually untrusted control planes
 - Path control by senders and receivers
 - Simpler routers (no forwarding tables)
 - Root of trust selectable by each ISD
- An **isolation architecture** for the **control plane**, but a **transparency architecture** for the **data plane**.



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ISP Deployment (Core AS)

- Core AS duties
 - Manage and distribute the ISD's TRC
 - Sign TRCs of neighboring ISDs and endorse other ISDs
 - Maintain a list of all recognized ISDs
 - Issue certificates to all ASes in the ISD
 - Provide connectivity to neighboring ISDs
 - Generate and disseminate inter-ISD path-segment construction beacons (PCBs), also called core PCBs
 - Generate and disseminate intra-ISD PCBs
 - Provide highly available services: beacon, name (RAINS), path, certificate, SIBRA, and time servers

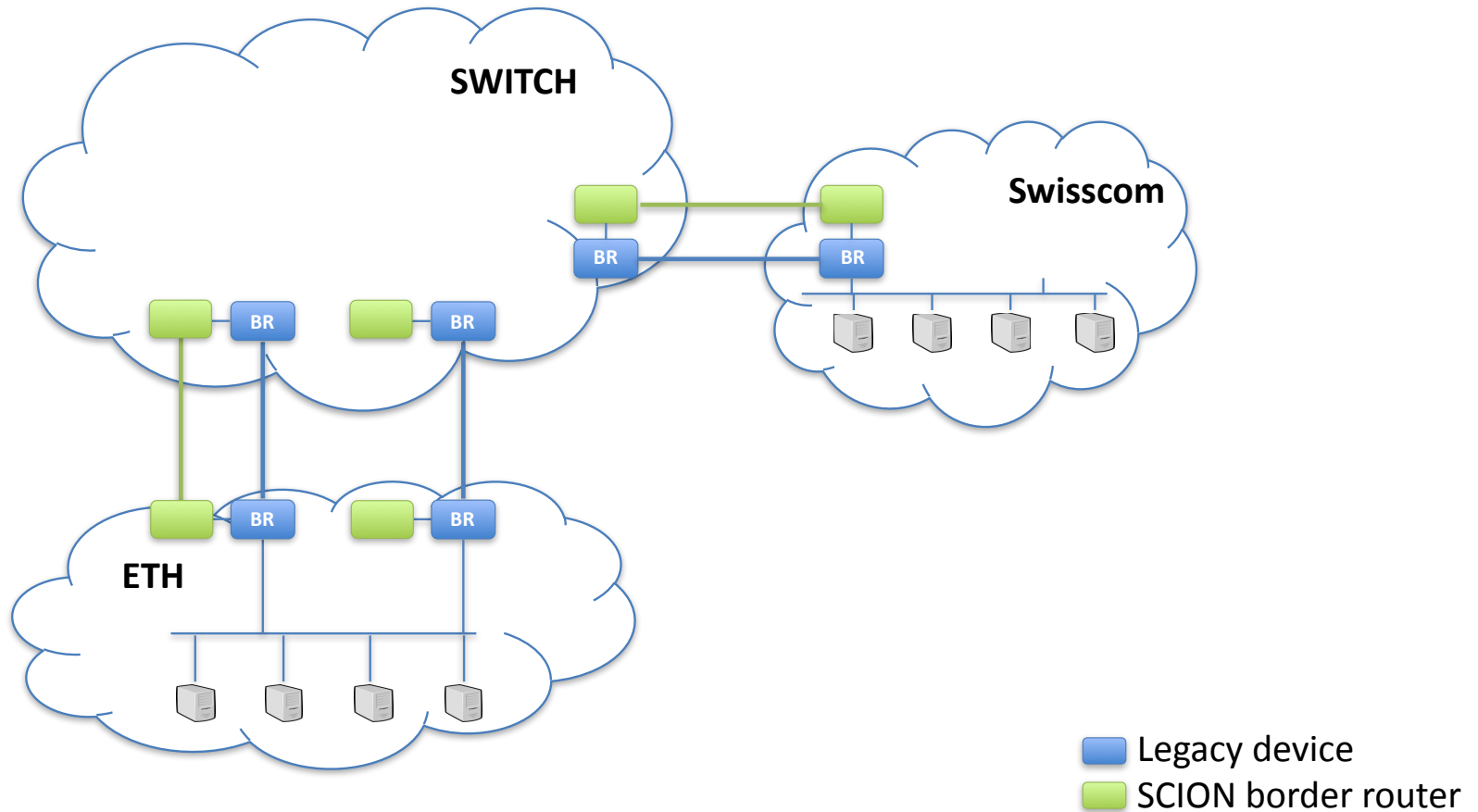
ISP Deployment (~~Core AS~~)

- ~~Core AS duties~~
 - ~~Manage and distribute the ISD's TRC~~
 - ~~Sign TRCs of neighboring ISDs and endorse other ISDs~~
 - ~~Maintain a list of all recognized ISDs~~
 - ~~Issue certificates to all ASes in the ISD~~
 - Provide connectivity to neighboring ~~ISDs~~ ASes
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 - ~~Generate and disseminate intra ISD PCBs~~
 - Provide ~~highly~~ available services: beacon, name (RAINS), path, certificate, SIBRA, and ~~time servers~~

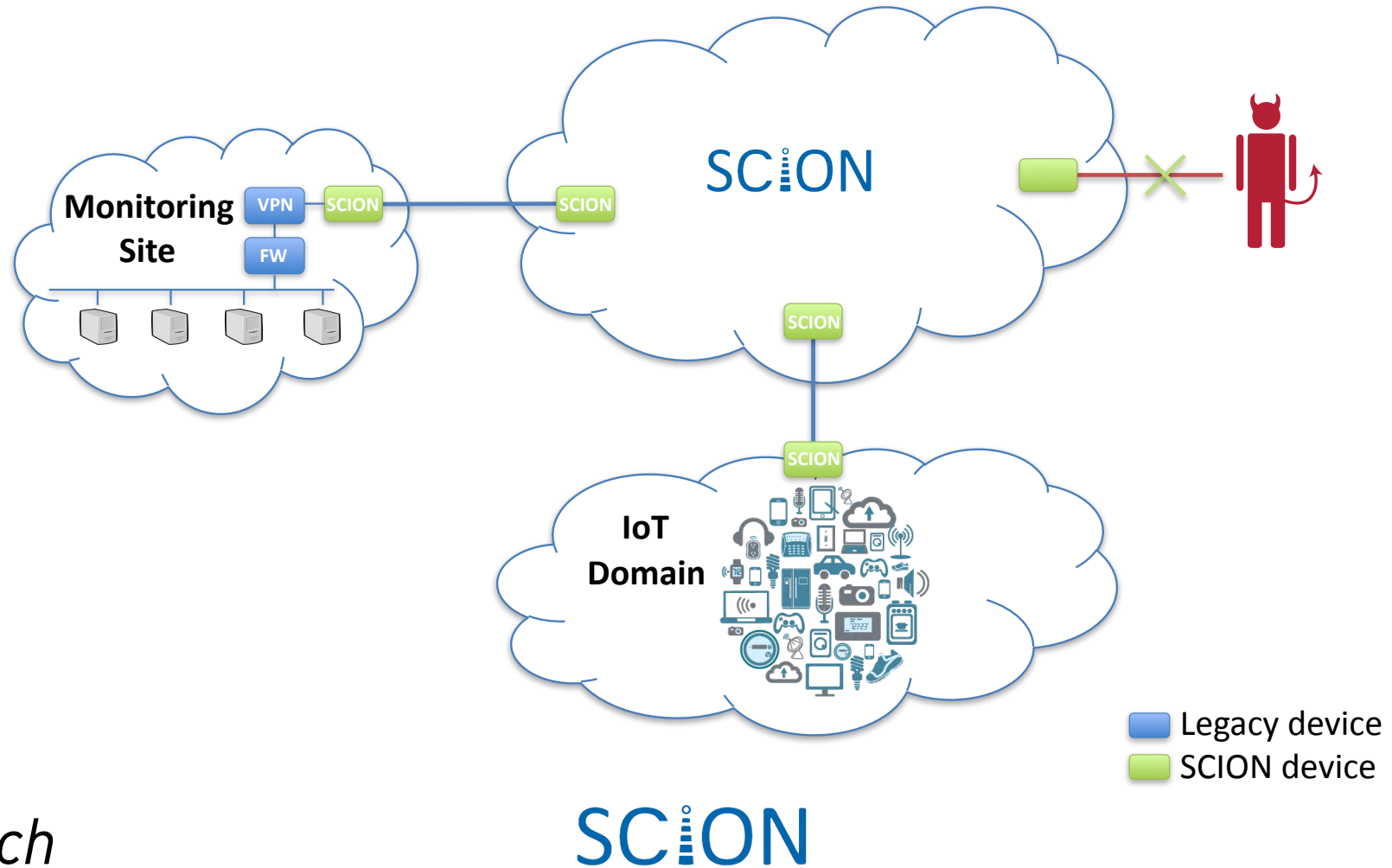
Leaf AS Deployment

- Leaf AS deployment tasks
 - Obtain AS certificate from core AS
 - Deploy servers: beacon, name (RAINS), path, certificate, SIBRA
- One single legacy PC suffices, e.g., attached to border router

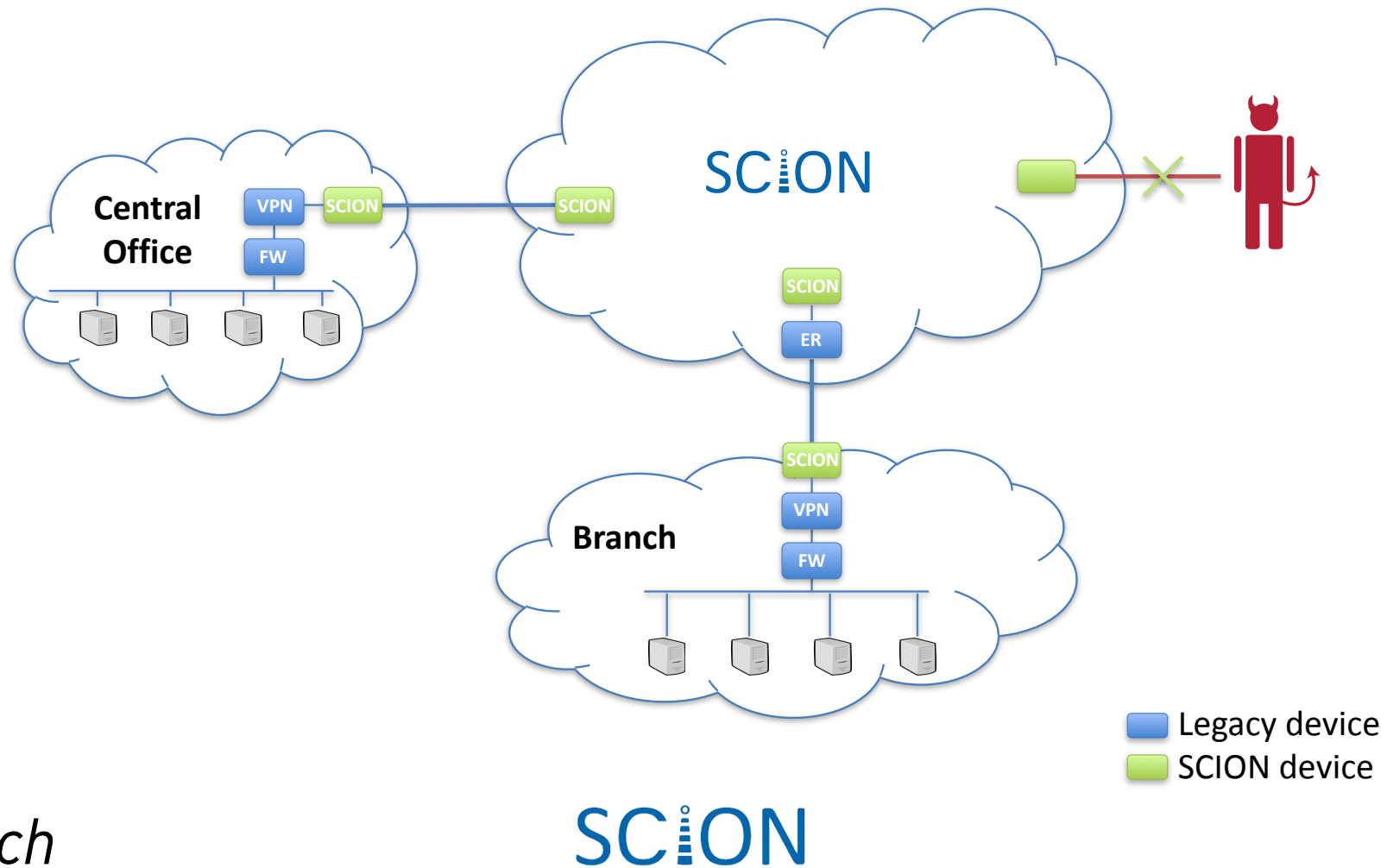
Deployment @ ETH



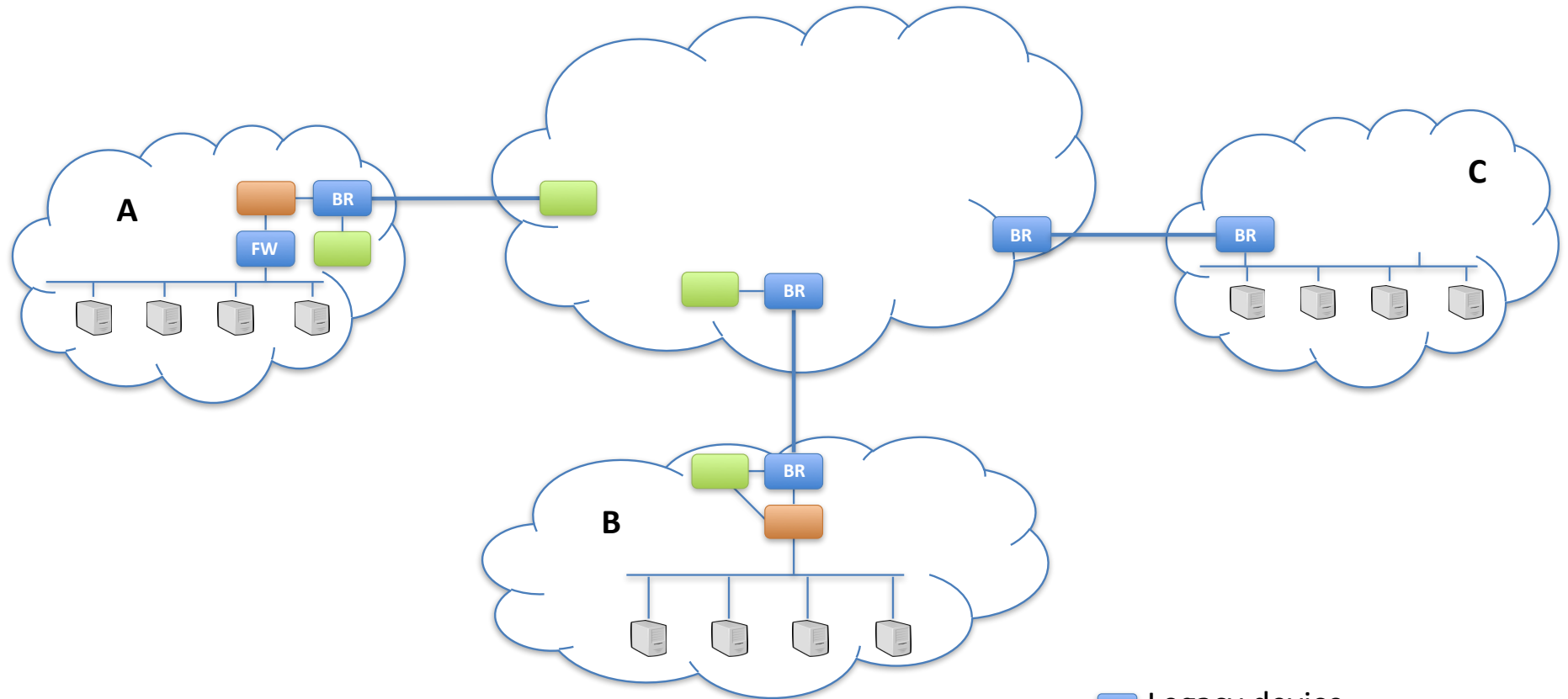
Use Case: IoT Protection through Default Off



Use Case: VPN-based Deployment

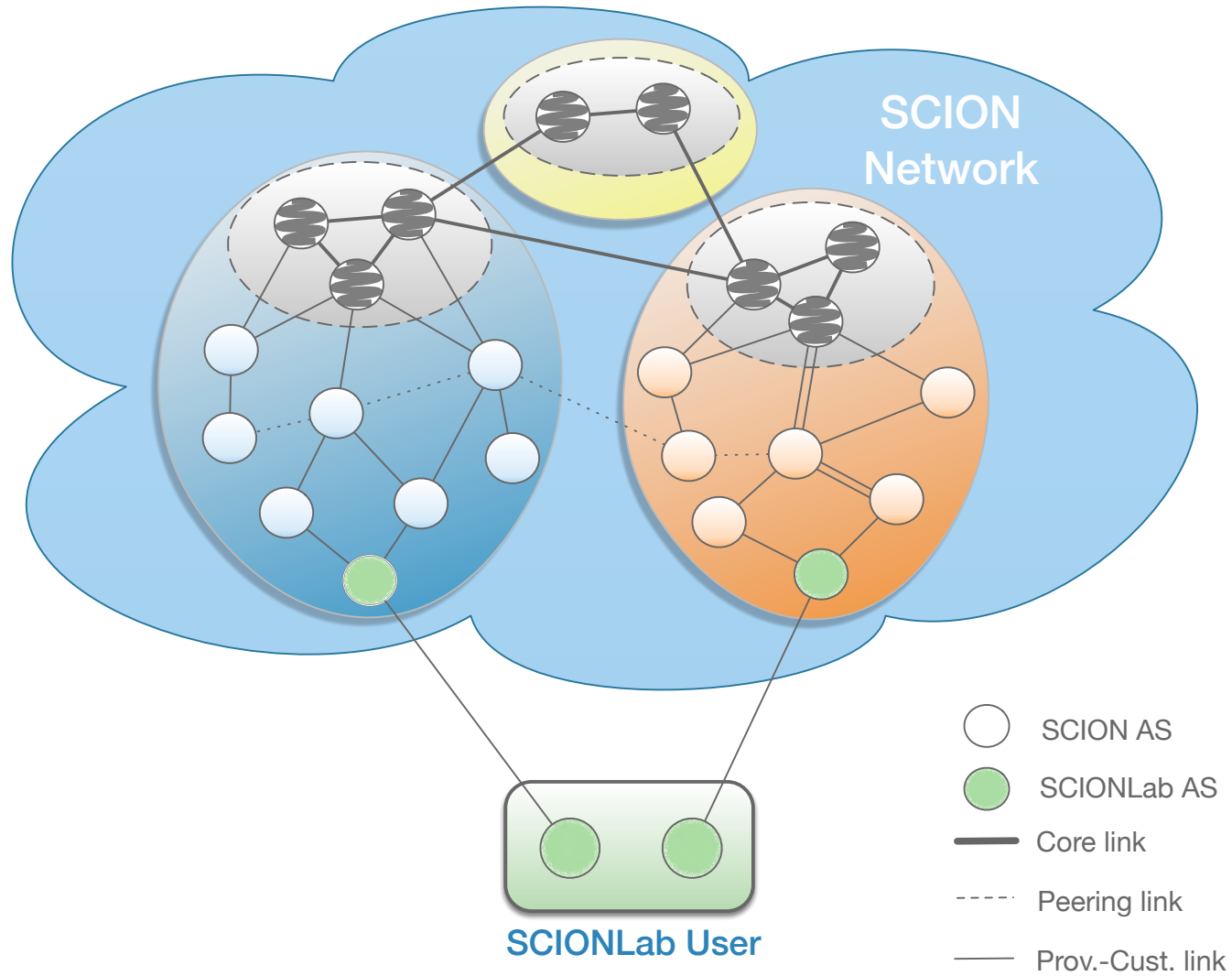


SCION-IP Gateway (SIG) Deployment

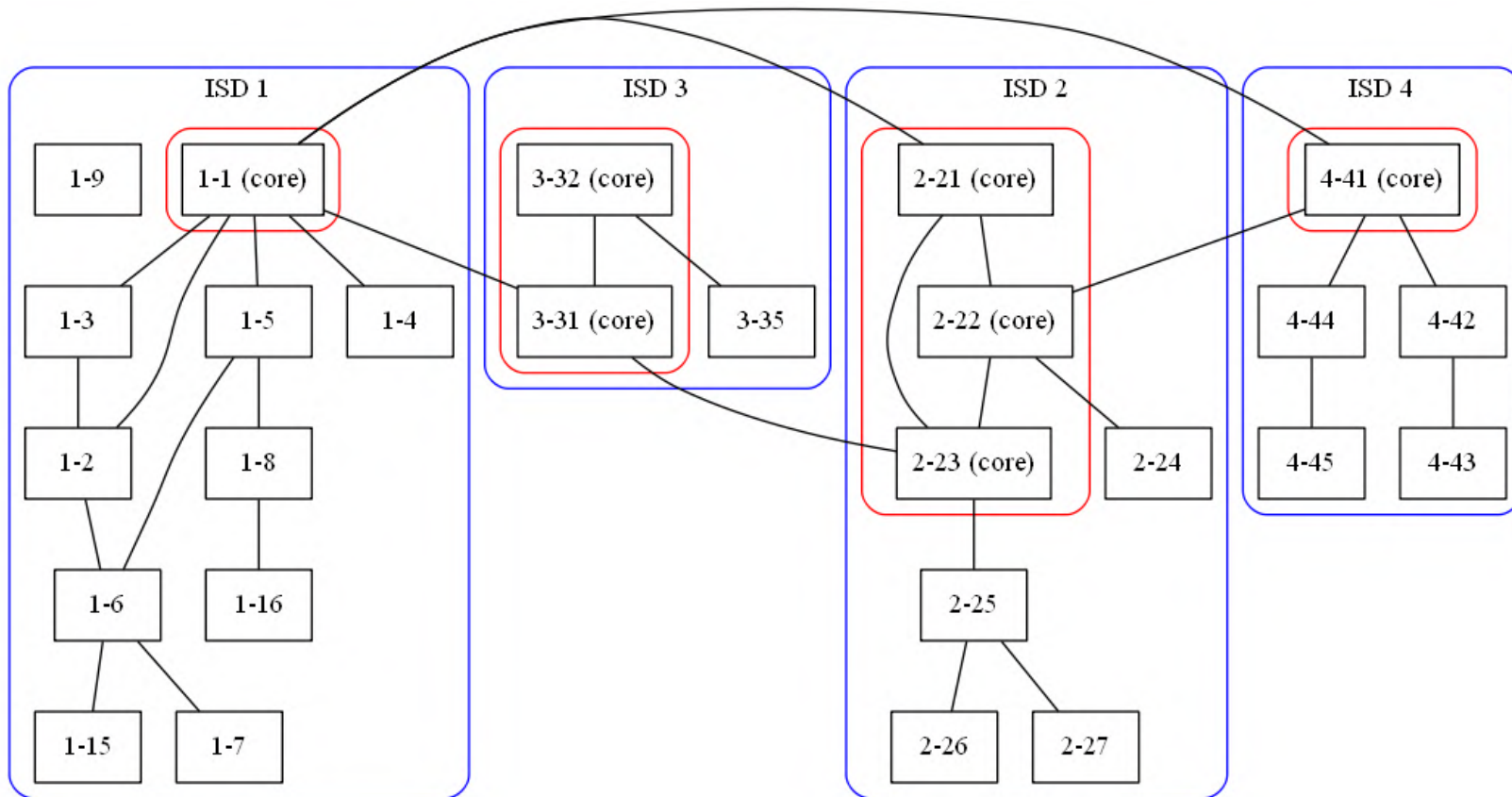


- Legacy device
- SCION border router
- SIG

SCIONLab



SCIONLab Network



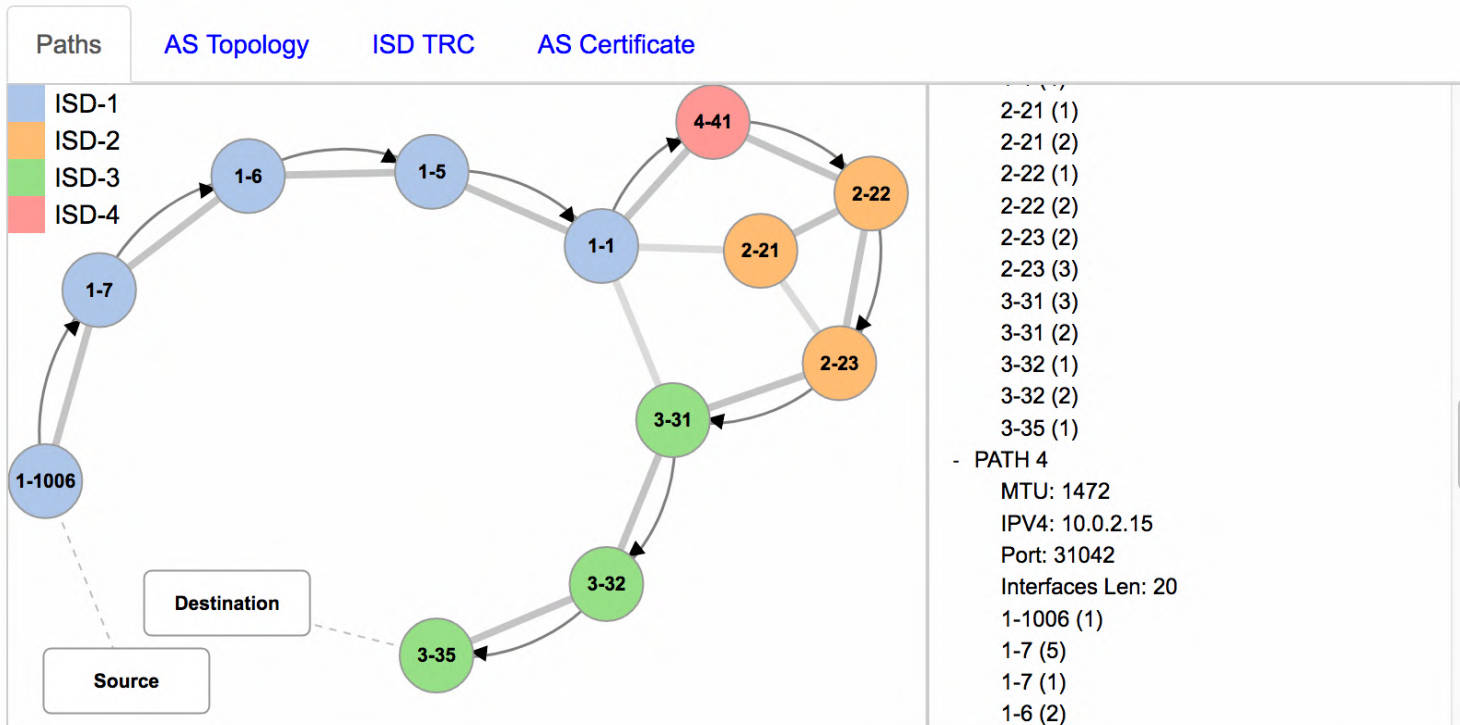
SCION Visualization System

SCION AS Visualization

- [SCION Website](#)
- [SCION on Github](#)
- [SCION Visualizations on Github](#)

Source AS: Destination AS: Data:

SCIOND IP Address:



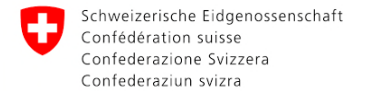
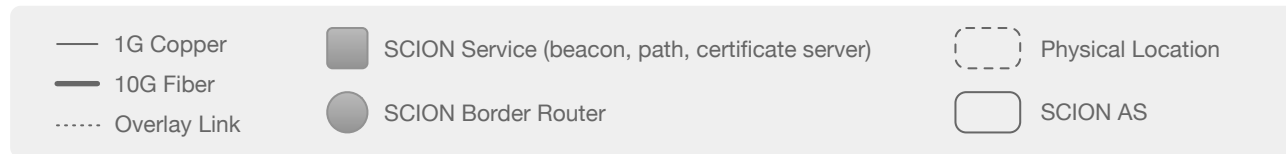
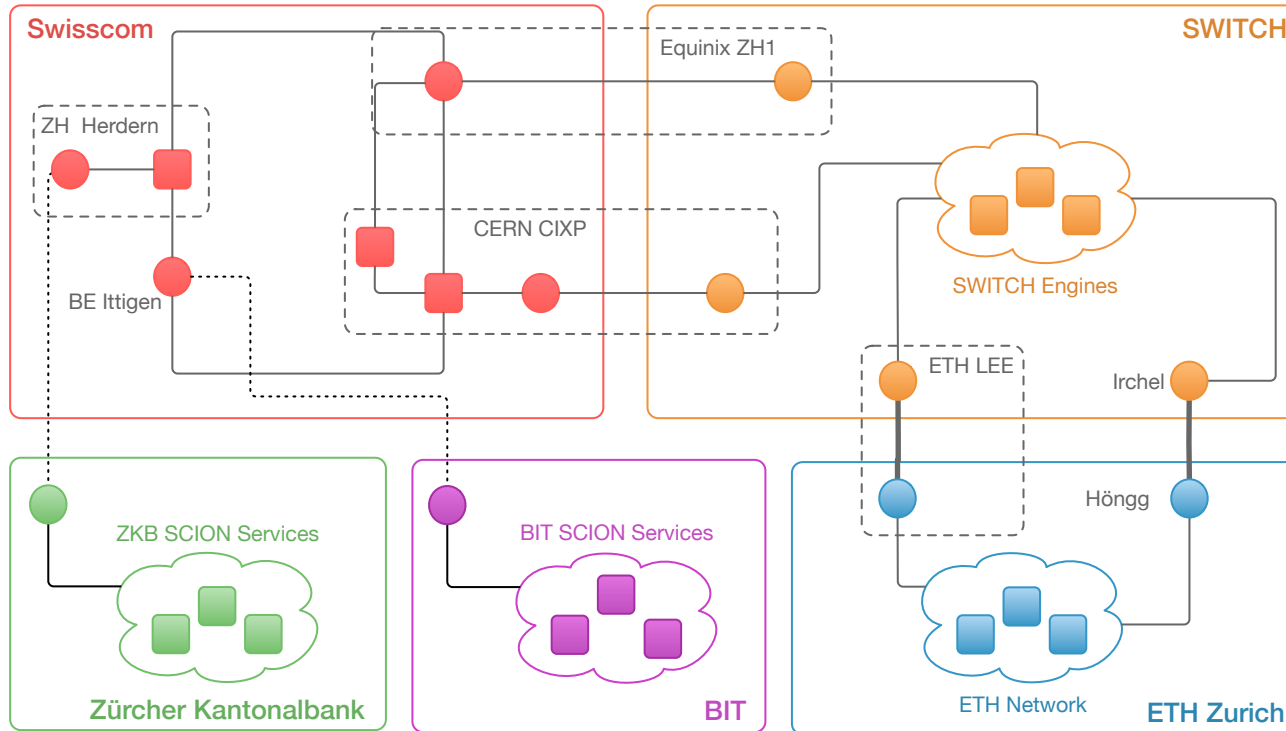
Application: IoT Access



AS Router Monitoring with Prometheus



Swiss SCION Network

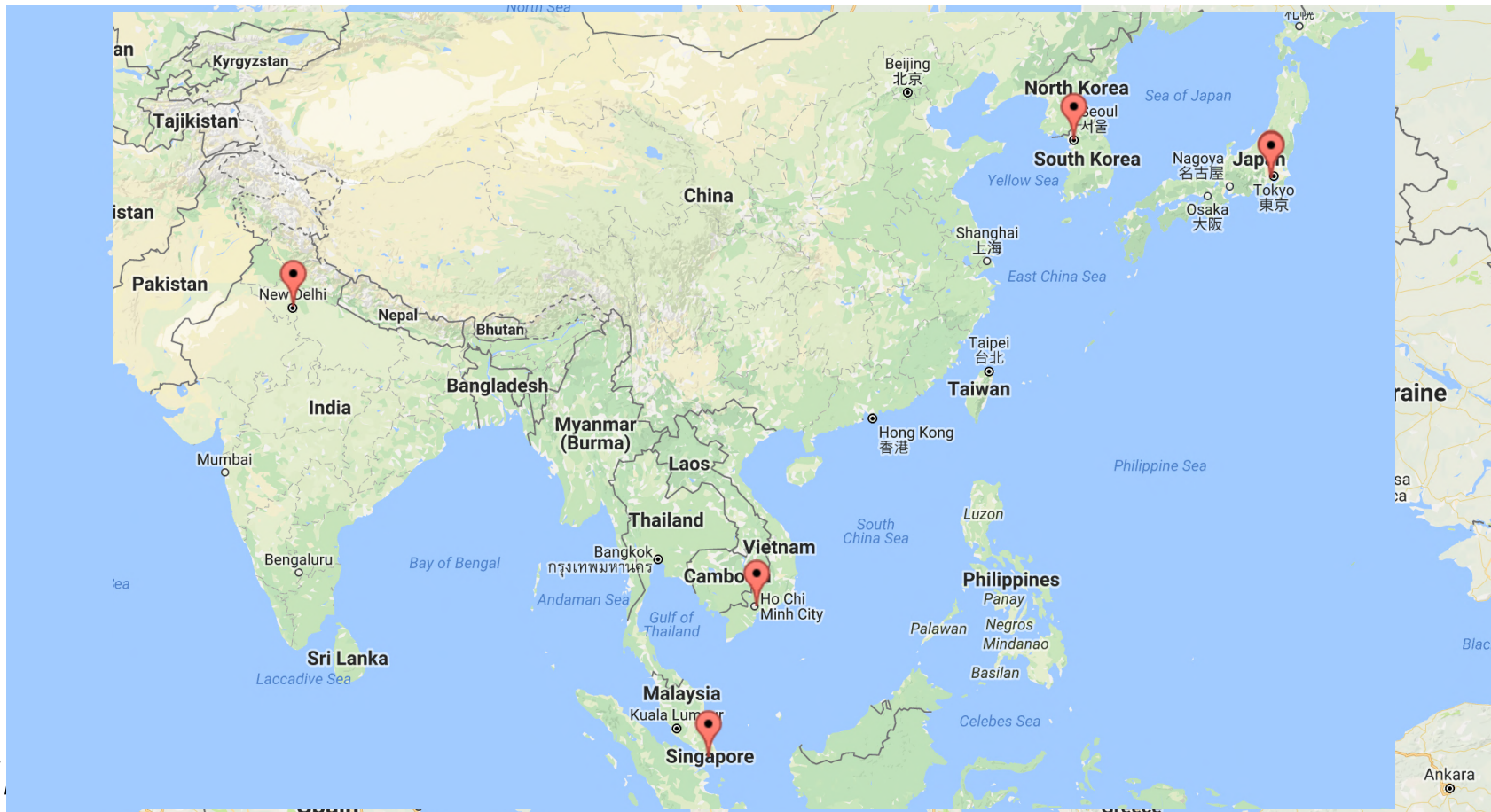


SCHWEIZERISCHE NATIONALBANK
 BANQUE NATIONALE SUISSE
 BANCA NAZIONALE SVIZZERA
 BANCA NAZIUNALA SVIZRA
 SWISS NATIONAL BANK



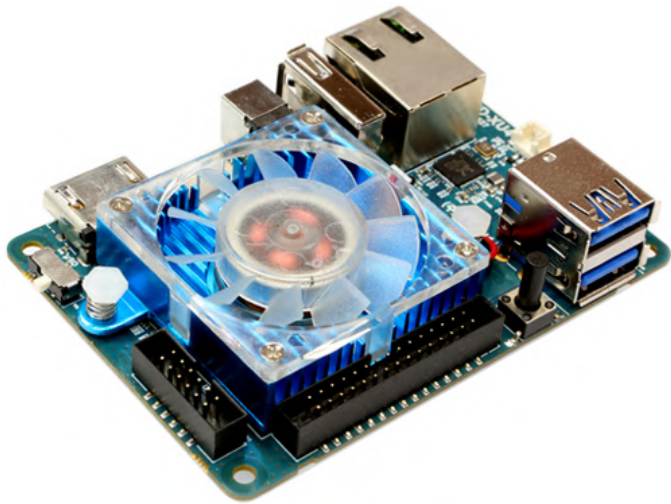
Growing Global Testbed

- Over 40 deployed SCION routers and servers

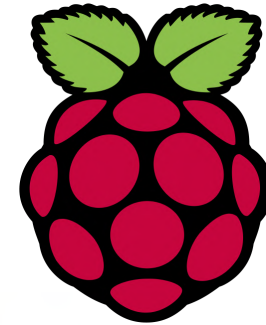


SCION AS runs on ODROID and Raspberry Pi

ODROID
Hardkernel



Raspberry Pi



Demos

- Fast failover and multipath demo
- Path control and geofencing demo

The screenshot shows a Grafana dashboard for a 'Border Router'. It features four line graphs: 'Bits per second - Received', 'Bits per second - Sent', 'Packets per second - Received', and 'Packets per second - Sent'. Below the graphs is a 'Configuration' section for ISD whitelists and a 'Topology' section with a network diagram. The diagram shows a central 'SCION' node connected to 'ETHz' and 'ZMB' nodes. A text box explains: 'Packet stream is received over Path 1, because it is the shortest path and has the lowest latency.'

All Paths	Path 1	Path 2	Path 3	Path 4
Packets Sent	2020	0	0	0
Packets Received	19	1	1	1
RTT (ms)	3.85	14.85	16.31	16.43
Loss Rates	0.00	0.00	0.00	0.00
Interface 1	1-2 (0)	1-3 (0)	1-3 (0)	1-3 (0)
Interface 2	1-2 (0)	1-2 (0)	1-2 (0)	1-2 (0)
Interface 3	1-2 (0)	1-2 (0)	1-2 (0)	1-2 (0)
Interface 4	1-2 (0)	1-2 (0)	1-2 (0)	1-2 (0)
Interface 5	1-2 (0)	1-2 (0)	1-2 (0)	1-2 (0)
Interface 6	1-2 (0)	1-2 (0)	1-2 (0)	1-2 (0)
Interface 7	1-2 (0)	1-2 (0)	1-2 (0)	1-2 (0)
Interface 8	1-2 (0)	1-2 (0)	1-2 (0)	1-2 (0)

The screenshot shows the SCION website with the tagline 'SCALABILITY, CONTROL, AND ISOLATION ON NEXT-GENERATION NETWORKS'. To the right is a network topology diagram with nodes labeled CH, DE, and FR. A legend identifies ISD-2 (orange), ISD-3 (green), ISD-4 (red), and ISD-5 (purple). Below the diagram is a 'Socket Proxy Data' table.

All Paths	Path 1	Path 2	Path 3	Path 4
Packets Sent	1	4	1	1
Packets Received	0	7	0	0
RTT (ms)	13.93	9.12	18.96	13.92
Loss Rates	0.00	0.00	0.00	0.00
Interface 1	1-2 (0)	1-2 (0)	1-2 (0)	1-2 (0)
Interface 2	1-2 (0)	1-2 (0)	1-2 (0)	1-2 (0)
Interface 3	1-2 (0)	1-2 (0)	1-2 (0)	1-2 (0)
Interface 4	1-1 (0)	1-1 (0)	1-1 (0)	1-1 (0)
Interface 5	1-1 (0)	1-1 (0)	1-1 (0)	1-1 (0)
Interface 6	1-1 (0)	1-1 (0)	1-1 (0)	1-1 (0)
Interface 7	1-1 (0)	1-1 (0)	1-1 (0)	1-1 (0)
Interface 8	1-1 (0)	1-1 (0)	1-1 (0)	1-1 (0)

Multiple paths are being used to load the webpage - a direct path between France and Switzerland and multiple paths through Germany.

Belief that Internet is Immutable

- Evidence appears overwhelming that Internet is immutable: IPv6, BGPSEC, DNSSEC, etc.
- However, benefits are limited, esp. for early deployers
- Our goal: provide many benefits, even for early adopters, such that one cannot turn back



Conclusions

- SCION is a secure Internet architecture that we can start using today
- Open source
- Numerous opportunities for researchers
 - Multipath routing architecture offers multitude of path choices for meaningful diverse path selection
 - Security: routing, DDoS, source authentication
 - Next-generation PKI architecture
- Natural quality scalability with increasing global adoption

SCION Projekt Team

- Netsec: Daniele Asoni, Chen Chen, Laurent Chuat, Sergiu Costea, Sam Hitz, Tobias Klausmann, Tae-Ho Lee, Chris Pappas, **Adrian Perrig**, Benjamin Rotenberger, Stephen Shirley, Jean-Pierre Smith, Pawel Szalachowski, Brian Trammell, Ercan Ucan
- Infsec: **David Basin**, Tobias Klenze, Christoph Sprenger, Thilo Weghorn
- Programming Methodology: Marco Eilers, **Peter Müller**





www.anapaya.net

Additional Information

- <https://www.scion-architecture.net>
 - Book
 - Papers
 - Videos
 - Newsletter signup
- <https://www.anapaya.net>
 - Commercializing SCION equipment
- <https://github.com/netsec-ethz/scion>